

# Size Ramsey numbers of forests versus double stars and brooms

Yuan Si

Qinghai Normal University, Xining, China; and  
Georgia State University, USA

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# Outline

## 1 Introduction

- Size Ramsey number
- Double stars
- Brooms

## 2 Main results

- Matchings versus double stars and brooms
- $P_3$  versus double stars and brooms

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# Basic definition

- We assume all graphs are finite, undirected and simple. For a graph  $G$ , the vertex set, edge set, order and size of  $G$  are denoted by  $V(G)$ ,  $E(G)$ ,  $v(G)$  and  $e(G)$ , respectively. The *neighborhood*,  $N_G(v)$  or  $N(v)$  for short, of a vertex  $v$  of  $G$  is the set of vertices adjacent to it. The *degree*,  $\deg_G(v)$  or  $\deg(v)$  for short, of a vertex  $v$  of  $G$  is the number of edges incident to it. The *minimum degree* of  $G$ ,  $\delta(G)$ , is the smallest of the degrees of vertices in  $G$  and the *maximum degree*,  $\Delta(G)$ , of  $G$  is the largest of the degrees of the vertices in  $G$ . For any subset  $X$  of  $V(G)$ , let  $G[X]$  denote the subgraph induced by  $X$ ; similarly, for any subset  $F$  of  $E(G)$ , let  $G[F]$  denote the subgraph induced by  $F$ . We use  $G - X$  to denote the subgraph of  $G$  obtained by removing all the vertices of  $X$  together with the edges incident with them from  $G$ . If  $X = \{v\}$ , we simply write  $G - v$  for  $G - \{v\}$ .

# Size Ramsey number

- If  $G$  and  $H$  are two graphs, we write  $F \rightarrow (G, H)$  to denote that  $G$  or  $H$  is a monochromatic subgraph of  $F$  in every 2-coloring of the edges of  $F$ . The *size Ramsey number*  $\widehat{r}(G, H)$  of two graphs  $G$  and  $H$ , introduced by Erdős et al., is defined as

$$\widehat{r}(G, H) = \min\{e(F) : F \rightarrow (G, H)\}.$$

# Basic proof idea

- Upper bound: Construct a graph  $F$  such that  $F \rightarrow (G, H)$ , then  $e(F)$  is the upper bound of  $\widehat{r}(G, H)$ .
- Lower bound: Any graph with a given number of edges can find a red/blue edge-coloring such that these graphs contain neither a red copy of  $G$  nor a blue copy of  $H$ .

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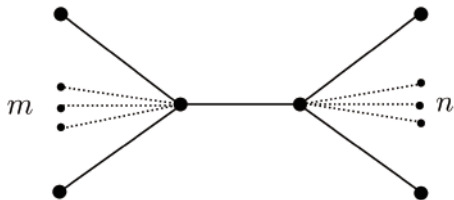
## 2 Main results

- Matchings versus double stars and brooms
- $P_3$  versus double stars and brooms



$D(m, n)$ 

A *double star* is a graph with two stars and an edge connecting their centers.



$$v(D(m, n)) = n+m+2, e(D(m, n)) = n+m+1, \Delta(D(m, n)) = \max\{n+1, m+1\}.$$

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## 1 Introduction

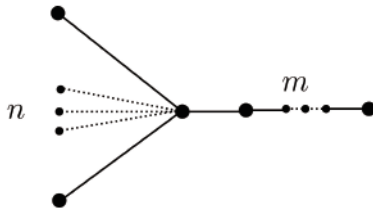
- Size Ramsey number
- Double stars
- **Brooms**

## 2 Main results

- Matchings versus double stars and brooms
- $P_3$  versus double stars and brooms

$B(m, n)$ 

A *broom* is a graph which is a path with a star at one end.



$$v(B(m, n)) = n + m + 1, e(B(m, n)) = n + m, \Delta(B(m, n)) = n + 1.$$

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# The subgraph induced by the red edges

## Fact 1

*To avoid red  $2P_2$ , the subgraph induced by the red edges in  $F$  is a star or a triangle.*

$\widehat{r}(sP_2, D(m, n))$ 

## Theorem 2.1

For  $n \geq m \geq 2$  and  $s \geq 2$ ,

$$s(n+1) + m \leq \widehat{r}(sP_2, D(m, n)) \leq s(n+m+1).$$

Upper bound of  $\widehat{r}(sP_2, D(m, n))$ 

Construct a graph  $F$  such that  $F \rightarrow (sP_2, D(m, n))$ :

$$sD(m, n)$$

$$e(sD(m, n)) = s(n + m + 1).$$



Lower bound of  $\widehat{r}(sP_2, D(m, n))$ 

Let  $F$  be a graph of size at most  $s(n+1) + m - 1$  and containing a copy of  $D(m, n)$ . Next use induction on  $s$ :

Assume  $s = 2$ .

- $\Delta(F) \leq n$ : no  $D(m, n)$  copy.
- $\Delta(F) \geq n + 1$ : Take a vertex  $v$  of maximum degree, and color the edges incident to  $v$  red and the other edges blue.

Suppose the conclusion holds for  $s - 1$ . Again, take a vertex  $v$  of maximum degree. Denote  $F' = F - v$ . Then

$$e(F') \leq s(n+1) + m - 1 - (n+1) = (s-1)(n+1) + m - 1.$$

Then by inductive hypothesis, there is a red-blue coloring of  $E(F')$  such that there is neither a red copy of  $(s-1)P_2$  nor a blue copy of  $D(m, n)$ . Now keep this coloring, and further color the edges red incident to  $v$ , then we get the red-blue coloring which we want.

$\widehat{r}(2P_2, D(1, n))$ 

## Theorem 2.2

*For  $n \geq 3$ ,*

$$\widehat{r}(2P_2, D(1, n)) = 2n + 4.$$

Upper bound of  $\widehat{r}(2P_2, D(1, n))$ 

By Theorem 2.1, we have  $\widehat{r}(2P_2, D(1, n)) \leq 2n + 4$ .  
Next we show  $\widehat{r}(2P_2, D(1, n)) \geq 2n + 4$ .

Lower bound of  $\widehat{r}(2P_2, D(1, n))$ 

Let  $F$  be a graph of size  $2n + 3$  and containing a copy of  $D(1, n)$ .  
Take  $w \in V(F)$  such that  $\deg(w) = \Delta(F)$ .

Let  $F' = F - w$ .

- $\Delta(F) \leq n$ : no  $D(1, n)$  copy.
- $\Delta(F) \geq n + 2$ :  $e(F') < e(D(1, n))$ , thus no  $D(1, n)$  copy.
- $\Delta(F) = n + 1$ : need further discussion.

# Lower bound of $\widehat{r}(2P_2, D(1, n))$

Thus assume that  $\Delta(F) = n + 1$ .

Then  $e(F') = n + 2 = e(D(1, n))$ .

- $F' \not\cong D(1, n)$ : we have done.
- $F' \cong D(1, n)$ : need further discussion.

Lower bound of  $\widehat{r}(2P_2, D(1, n))$ 

Assume  $F' \cong D(1, n)$ .

Let  $u$  and  $v$  be the  $n$ -center and 1-center, respectively.

- $N(w) \cap V(F') = \emptyset$ : color one edge of  $F'$  red and the other edges blue.

Thus assume  $N(w) \cap V(F') \neq \emptyset$ .

- $v \in N(w) \cap V(F')$ : color the edges incident to  $v$  red and color the other edges blue.

Lower bound of  $\widehat{r}(2P_2, D(1, n))$ 

Thus assume  $v \notin N(w) \cap V(F')$ .

- $u \in N(w) \cap V(F')$ : color the edge  $uw$  red and color the other edges blue.

Thus assume  $u \notin N(w) \cap V(F')$ .

- $\exists x \in N_{F'}(u)$  such that  $x \in N(w) \cap V(F')$ : color  $ux$  and  $wx$  red and color the other edges blue.

Thus assume no neighbor of  $u$  belong to  $N(w) \cap V(F')$ .

- Denote by  $v_1$  the other neighbor of  $v$  than  $u$  in  $F'$  such that  $N(w) \cap V(F') = \{v_1\}$ : color  $vv_1$  and  $wv_1$  red and other edges blue.

And such coloring is what we want.

$\widehat{r}(sP_2, D(1, n))$ 

Next by an induction on  $s$ , we have the following result:

## Lemma 2

Assume  $F$  is a graph of size  $s(n+2) - 1$  where  $s \geq 2$  and  $n \geq \lceil (s^2 + 3s - 2)/2 \rceil$ . Denote by  $W$  the set of vertices of degree at least  $n+1$  in  $F$ . If  $|W| \geq s$ , then  $|W| = s$ .

## Theorem 2.3

For  $s \geq 2$  and  $n \geq \lceil (s^2 + 3s - 2)/2 \rceil$ ,

$$\widehat{r}(sP_2, D(1, n)) = s(n+2).$$

Skipping the proof.



$\widehat{r}(2P_2, D(2, 2))$ 

## Theorem 2.4

$$\widehat{r}(2P_2, D(2, 2)) = 10.$$

Skipping the proof.

Problem of  $\widehat{r}(2P_2, D(n, n))$ **Problem 2.1**

For  $n \geq 3$ ,

$$\widehat{r}(2P_2, D(n, n)) = ?.$$

$\widehat{r}(2P_2, B(m, n))$ 

## Theorem 2.5

For  $m \geq 4, n \geq 1$ ,

$$\widehat{r}(2P_2, B(m, n)) \leq 2n + 2m - 2.$$

For  $m \geq 3, n \geq m + 2$ ,

$$\widehat{r}(2P_2, B(m, n)) \geq 2n + m + 2.$$

Upper bound of  $\widehat{r}(2P_2, B(m, n))$ 

Construct a graph  $F$  such that  $F \rightarrow (2P_2, B(m, n))$ :

let  $F$  be a graph obtained from a  $2m$ -cycle  $v_1 v_2 \dots v_{2m}$  and two stars  $K_{1, n-1}, K_{1, n-1}$  with centers  $u, v$ , respectively, by identifying  $u_1, v_1$  and  $u_2, v_m$  (that is  $u_1 = v_1$  and  $u_2 = v_m$ ), which show in Figure 1(a).

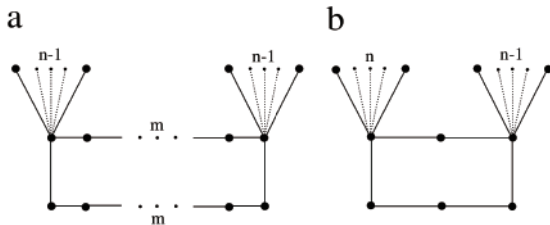


Figure 1: Graphs for Theorem 2.5 and Corollary 2.1.

Lower bound of  $\widehat{r}(2P_2, B(m, n))$ 

Let  $F$  be a graph with at most  $2n + m + 1$ .

We can assume that  $e(F) = 2n + m + 1$ .

Take  $w \in V(F)$  such that  $\deg(w) = \Delta(F)$ .

Let  $F' = F - w$ .

- $\Delta(F) \leq n$ : no  $B(m, n)$  copy.
- $\Delta(F) \geq n + 2$ :  $e(F') < e(B(m, n))$ , thus no  $B(m, n)$  copy.
- $\Delta(F) = n + 1$ : need further discussion.

Lower bound of  $\widehat{r}(2P_2, B(m, n))$ 

Suppose  $\Delta(F) = n + 1$ . Let  $s$  be the number of vertices of degree  $n + 1$ .

## Claim 3

$$s \leq 2.$$

## Proof.

Assume, to the contrary, that  $s \geq 3$ . Then there exist three vertices, say  $u_1, u_2, u_3$ , such that  $\deg_F(u_i) = \Delta(F) = n + 1$  for  $i = 1, 2, 3$ , and hence there are at least  $3(n + 1) - 3 \geq 2n + m + 2 > e(F)$ , since  $n \geq m + 2$ , a contradiction.  $\blacksquare$

Lower bound of  $\widehat{r}(2P_2, B(m, n))$ 

Suppose  $s = 2$ . Then there exist two vertices  $u, v \in V(F)$  such that  $\deg_F(u) = \deg_F(v) = \Delta(F) = n + 1$ .

- $uv \in E(F)$ : color the edge  $uv$  red, and then color the other edges blue.
- $uv \notin E(F)$  and  $N_F(u) \cap N_F(v) \neq \emptyset$ , then  $\exists x \in N_F(u) \cap N_F(v)$ : color the edges  $ux, vx$  red, and then color the other edges blue.
- $uv \notin E(F)$  and  $N_F(u) \cap N_F(v) = \emptyset$ : note that  $e(F - u - v) = m - 1$  and  $F - v \cong B(m, n)$ ,  $F - u \cong B(m, n)$ , then these  $m - 1$  edges must form a path. We color one edge from this path red, and then color the other edges blue.

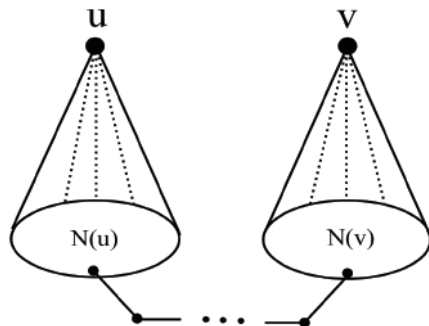
Lower bound of  $\widehat{r}(2P_2, B(m, n))$ 

Figure 2: Graph for the Proof of Theorem 2.5.



$\widehat{r}(2P_2, B(3, n))$  and  $\widehat{r}(2P_2, B(4, n))$ 

## Corollary 2.1

For  $n \geq 5$ ,

$$\widehat{r}(2P_2, B(3, n)) = 2n + 5.$$

## Corollary 2.2

For  $n \geq 6$ ,

$$\widehat{r}(2P_2, B(4, n)) = 2n + 6.$$

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# The subgraph induced by the red edges

## Fact 4

*To avoid red  $P_3$ , the subgraph induced by the red edges in  $F$  is a matching.*

$\widehat{r}(P_3, D(m, n))$ 

## Theorem 2.6

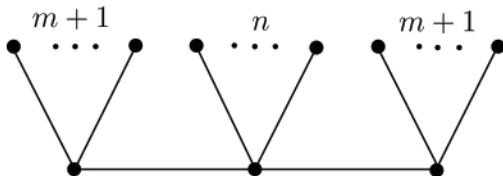
For  $n \geq m \geq 2$ ,

$$n + 2m + 1 \leq \widehat{r}(P_3, D(m, n)) \leq n + 2m + 4.$$

Upper bound of  $\widehat{r}(P_3, D(m, n))$ 

Construct a graph  $F$  such that  $F \rightarrow (P_3, D(m, n))$ :

Let  $uvw$  be a path of order 3, and  $F$  be a graph obtained from three stars  $K_{1,m+1}$ ,  $K_{1,n}$  and  $K_{1,m+1}$  by identifying the center of  $K_{1,m+1}$  and  $u$ , the center of  $K_{1,n}$  and  $v$  and the center of  $K_{1,m+1}$  and  $w$ , respectively.



Skipping the proof of lower bound.

# Remark

From Theorem 2.6,  $3n + 1 \leq \widehat{r}(P_3, D(n, n)) \leq 3n + 4$  if we take  $m = n$ . Furthermore, we will show  $\widehat{r}(P_3, D(n, n)) = 3n + 4$  for  $n \geq 5$ . For a maximum matching  $M$  of a graph  $F$ , we denote

$$V(M, n) = \{v \in V(F) \mid \deg(v) \geq n + 1 \text{ and } v \text{ covered by } M\}.$$

Then we have the following lemma and theorem.

$\widehat{r}(P_3, D(n, n))$ 

## Lemma 5

Let  $F$  be a graph with  $e(F) \leq 3n + 3$  and  $\Delta(F) \geq n + 1$ , where  $n \geq 5$ . Let  $M$  be a maximum matching of  $F$  such that  $|V(M, n)|$  is maximized. For any two vertices  $u, v$  of degree at least  $n + 1$ , if  $uv \in E(F)$ , then  $M$  covers both  $u$  and  $v$ .

## Theorem 2.7

For  $n \geq 5$ ,

$$\widehat{r}(P_3, D(n, n)) = 3n + 4.$$

Skipping the proof.

$\widehat{r}(P_3, D(1, n))$ 

## Theorem 2.8

For  $n \geq 4$ ,

$$\widehat{r}(P_3, D(1, n)) = n + 5.$$

Skipping the proof.



$$\widehat{r}(P_3, B(3, n))$$

### Theorem 2.9

For  $n \geq 5$ ,

$$\widehat{r}(P_3, B(3, n)) = n + 6.$$

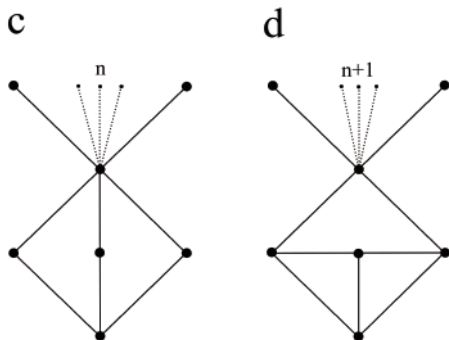
Upper bound of  $\widehat{r}(P_3, B(3, n))$  and  $\widehat{r}(P_3, B(4, n))$ 

Figure 3: Graphs for Theorem 2.9 and Theorem 2.10.

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

Next we show  $\widehat{r}(P_3, B(3, n)) \geq n + 6$ . Let  $F$  be a graph with at most  $n + 5$ . We can assume that  $e(F) = n + 5$ . Let  $s$  be the number of vertices of degree  $\geq n + 1$ .

## Claim 6

$$s = 1.$$

## Proof.

Assume, to the contrary, that  $s \geq 2$ . Then there exist two vertices, say  $u_1, u_2$ , such that  $\deg_F(u_i) \geq n + 1$  for  $i = 1, 2$ , and hence there are at least  $2(n + 1) - 1 \geq n + 6 > e(F)$ , since  $n \geq 5$ , a contradiction. ■

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

Assume  $\deg(v) = \Delta(F)$ .

- $\Delta(F) \leq n$ : no  $B(3, n)$  copy.
- $\Delta(F) = n + 1$ : choose an edge  $e$  incident to  $v$  in  $F$  and color it red, and the other edges in  $F - e$  are colored blue.
- $\Delta(F) \geq n + 3$ : color one edge which not incident  $v$  red, and the other edges blue.
- $\Delta(F) = n + 2$ : need further discussion.

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

Suppose  $\Delta(F) = n + 2$  and  $\deg(v) = \Delta(F)$ .

- $F - v \not\cong K_3$  and  $F - v \not\cong K_{1,3}$ : there is a  $2P_2$  in  $F - v$ , and we color  $2P_2$  red, and the other edges of  $F$  are colored blue..
- $F - v \cong K_3$ : need further discussion.
- $F - v \cong K_{1,3}$ : need further discussion.

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

Suppose  $F - v \cong K_3$  or  $F - v \cong K_{1,3}$ . Let  $N(v) = \{v_i \mid 1 \leq i \leq n + 2\}$ . The center vertex of  $K_{1,3}$  is denoted as  $u$ , and the other three degree vertices in  $K_{1,3}$  are denoted as  $u_1$ ,  $u_2$  and  $u_3$ .

Since  $F$  is connected,  $|N(v) \cap V(F - v)| \neq \emptyset$ .

If  $|N(v) \cap V(F - v)| = 1$ , let  $\{x\} = N(v) \cap V(F - v)$ , then we color the edge  $vx$  red, the other edges in  $F$  blue. One can easily check that there is neither a red copy of  $P_3$  nor a blue copy of  $B(3, n)$ .

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

## Case 1

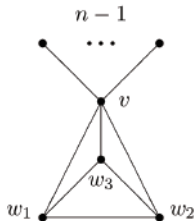
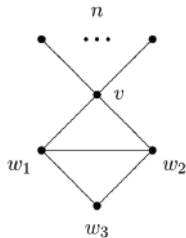
$$F - v \cong K_3$$

## Subcase 1.1

If  $|N(v) \cap V(F - v)| = 2$ , then we color the edge  $vv_i$  ( $v_i \in N(v)$  and  $\deg(v_i) = 1$ ) red, the other edges in  $F$  blue. Then  $v(F - v_i) = n + 3 < v(B(3, n))$ , so there is neither a red copy of  $P_3$  nor a blue copy of  $B(3, n)$ .

## Subcase 1.2

If  $|N(v) \cap V(F - v)| = 3$ , then  $v(F) = n + 3 < v(B(3, n))$ , which implies no  $B(3, n)$  copy.

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 



Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

## Case 2

$$F - v \cong K_{1,3}$$

## Subcase 2.1

If  $N(v) \cap V(F - v) = \{u, u_1\}$ , then we color the edges  $uu_2$  and  $vv_i$  ( $v_i \in N(v)$  and  $\deg(v_i) = 1$ ) red, the other edges in  $F$  blue. Then  $v(F - u_2 - v_i) = n + 3 < v(B(3, n))$ , so there is neither a red copy of  $P_3$  nor a blue copy of  $B(3, n)$ .

## Subcase 2.2

If  $N(v) \cap V(F - v) = \{u_1, u_2\}$ , then we color the edges  $uu_3$  and  $vv_i$  ( $v_i \in N(v)$  and  $\deg(v_i) = 1$ ) red, the other edges in  $F$  blue. Then  $v(F - u_3 - v_i) = n + 3 < v(B(3, n))$ , so there is neither a red copy of  $P_3$  nor a blue copy of  $B(3, n)$ .

Lower bound of  $\widehat{r}(P_3, B(3, n))$ 

## Subcase 2.3

*If  $|N(v) \cap V(F - v)| = 3$ , then we color the edge  $vv_i$  ( $v_i \in N(v)$  and  $\deg(v_i) = 1$ ) red, the other edges in  $F$  blue. Then  $v(F - v_i) = n + 3 < v(B(3, n))$ , so there is neither a red copy of  $P_3$  nor a blue copy of  $B(3, n)$ .*

## Subcase 2.4

*If  $|N(v) \cap V(F - v)| = 4$ , then  $v(F) = n + 3 < v(B(3, n))$ , which implies no  $B(3, n)$  copy.*

$\widehat{r}(P_3, B(4, n))$ 

## Theorem 2.10

*For  $n \geq 7$ ,*

$$\widehat{r}(P_3, B(4, n)) = n + 8.$$

Skipping the proof.

Problem of  $\widehat{r}(P_3, B(m, n))$ 

A good upper and lower bound.

**Problem 2.2**

For  $n \geq ?$ ,  $m \geq ?$ ,

$$? \leq \widehat{r}(P_3, B(m, n)) \leq ?.$$

# The End

Thank you for your  
attention !